## Barycentric Lagrange

Let's work on evaluating the Lagrange interpolator with knots in the xy-plane. Here we have n+1 nodes:  $x_0 < x_1 < \ldots < x_n$ , and n+1 knots:  $(x_i, y_i)$ , i = 0:n. The vector of y-values may just be data, or we may know that  $y_i = f(x_i)$  for some function f. The Lagrange interpolator of degree n (or less) for a function f on these knots gives

$$f = P + R$$
, where  $P = \sum_{i=0}^{n} y_j L_{nj}$ , and  $R$  is the truncation error,

and we want to find the number P(x), for some  $x \in \mathbb{R}$  (maybe many such values).

Each Lagrange basis function,  $L_{nj}$ , is a polynomial of degree n, and we have n+1 of them to evaluate, so it looks like about  $2n^2$  flops are required. We'll rewrite P in a way that cuts this to about  $n^2/2$  flops. We begin by noting that the denominator of  $L_{nj}$  is a scalar which only depends on the nodes:

$$L_{nj}(x) = \prod_{i=0, i \neq j}^{n} \frac{x - x_i}{x_j - x_i} \equiv w_j \prod_{i=0, i \neq j}^{n} (x - x_i), \quad \text{with} \quad w_j = \frac{1}{\prod_{i=0, i \neq j}^{n} (x_j - x_i)}.$$

If we let  $l(x) = \prod_{i=0}^{n} (x - x_i)$  (which doesn't depend on f), then

$$L_{nj}(x) = w_j \frac{l(x)}{x - x_j}$$
 and thus  $P(x) = \sum_{j=0}^n \frac{y_j w_j l(x)}{x - x_j}$ .

Noticing that l(x) is common to all terms, we can write

$$P(x) = l(x) \sum_{j=0}^{n} \frac{y_j w_j}{x - x_j}$$

While this is a good starting point for computing P(x), we can push it a bit further by using the Lagrange interpolator on these nodes for the function  $g(x) \equiv 1$ . The  $y_j$  for this gare  $y_j = 1$ , j = 0:n, and since g is a polynomial of degree  $\leq n$ , its interpolator is g itself (there is no remainder). This gives the identity

$$l(x)\sum_{j=0}^{n}\frac{w_j}{x-x_j} = 1 \qquad \text{for all } x.$$

Dividing P(x) by 1 (and cancelling the l(x) factors) gives the *Barycentric form* of the Lagrange interpolator:

$$P(x) = \frac{\sum_{j=0}^{n} \frac{y_j w_j}{x - x_j}}{\sum_{j=0}^{n} \frac{w_j}{x - x_j}}.$$

This form requires computing the  $w_j$ 's (which only depend on the knots, but be careful of over/under-flow), and can be done in about  $n^2/2$  flops. After that this Barycentric form only requires 5n + 1 flops per evaluation point. There is more to be said if you want to evaluate P for many points all at once... but that is another course (look up polynomial evaluation and the FFT and/or Cauchy matrices).